

Machine-Level Programming IV: Data

CSci 2021: Machine Architecture and Organization
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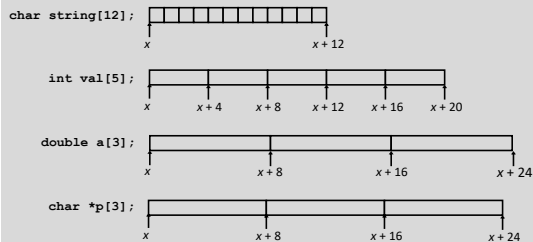
Based on slides originally by:
Randy Bryant, Dave O'Hallaron

Today

- Arrays
 - One-dimensional
 - Multi-dimensional (nested)
 - Multi-level
- Structures
 - Allocation
 - Access
 - Alignment

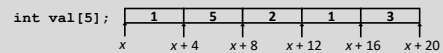
Array Allocation

- Basic Principle
 - $T\ A[L];$
 - Array of data type T and length L
 - Contiguously allocated region of $L * \text{sizeof}(T)$ bytes in memory



Array Access

- Basic Principle
 - $T\ A[L];$
 - Array of data type T and length L
 - Identifier A can be used as a pointer to array element 0: Type T^*

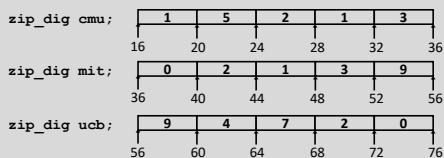


Reference	Type	Value
val[4]	int	3
val	int *	x
val+1	int *	x+4
&val[2]	int *	x+8
val[5]	int	??
*(val+1)	int	5
val + i	int *	x + 4i

Array Example

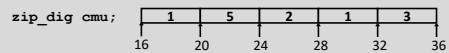
```
#define ZLEN 5
typedef int zip_dig[ZLEN];

zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```



- Declaration "zip_dig cmu" equivalent to "int cmu[5]"
- Example arrays were allocated in successive 20 byte blocks
 - Not guaranteed to happen in general

Array Accessing Example



```
int get_digit
(zip_dig z, int digit)
{
    return z[digit];
}
```

x86:

```
# %rdi = z
# %rsi = digit
movl (%rdi,%rsi,4), %eax # z[digit]
```

- Register `%rdi` contains starting address of array
- Register `%rsi` contains array index
- Desired digit at `%rdi + 4*%rsi`
- Use memory reference `(%rdi,%rsi,4)`

Array Loop Example

```
void zincr(zip_dig z) {
    size_t i;
    for (i = 0; i < ZLEN; i++)
        z[i]++;
}
```

```
# %rdi = z
movl $0, %eax      # i = 0
jmp .L3            # goto middle
.L4:                # loop:
addl $1, (%rdi,%rax,4) # z[i]++
addq $1, %rax      # i++
.L3:                # middle
cmpq $4, %rax      # i:4
jbe .L4            # if <=, goto loop
rep; ret
```

Multidimensional (Nested) Arrays

Declaration

- T $A[R][C]$;
- 2D array of data type T
- R rows, C columns
- Type T element requires K bytes

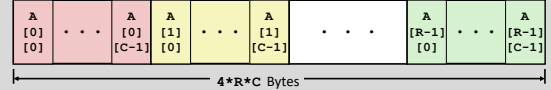
Array Size

- $R * C * K$ bytes

Arrangement

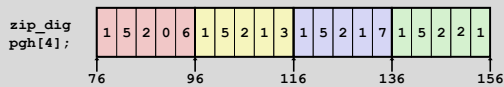
- Row-Major Ordering

```
int A[R][C];
```



Nested Array Example

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
{
    {1, 5, 2, 0, 6},
    {1, 5, 2, 1, 3},
    {1, 5, 2, 1, 7},
    {1, 5, 2, 2, 1}
};
```



"zip_dig pgh[4]" equivalent to "int pgh[4][5]"

- Variable `pgh`: array of 4 elements, allocated contiguously
- Each element is an array of 5 int's, allocated contiguously

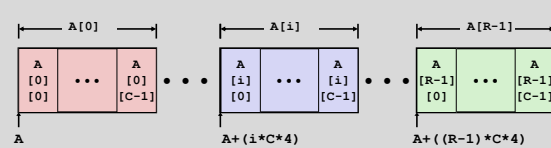
"Row-Major" ordering of all elements in memory

Nested Array Row Access

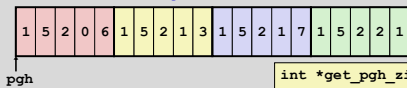
Row Vectors

- $A[i]$ is array of C elements
- Each element of type T requires K bytes
- Starting address $A + i * (C * K)$

```
int A[R][C];
```



Nested Array Row Access Code



```
int *get_pgh_zip(int index)
{
    return pgh[index];
}
```

```
# %rdi = index
leaq (%rdi,%rdi,4),%rax # 5 * index
leaq pgh(,%rax,4),%rax # pgh + (20 * index)
```

Row Vector

- `pgh[index]` is array of 5 int's
- Starting address `pgh + 20 * index`

Machine Code

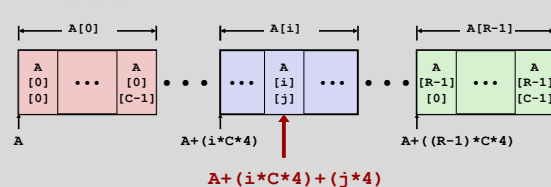
- Computes and returns address
- Compute as `pgh + 4 * (index * 4)`

Nested Array Element Access

Array Elements

- $A[i][j]$ is element of type T , which requires K bytes
- Address $A + i * (C * K) + j * K = A + (i * C + j) * K$

```
int A[R][C];
```



Nested Array Element Access Code

1	5	2	0	6	1	5	2	1	3	1	5	2	1	7	1	5	2	2	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

pgh

```
int get_pgh_digit
(int index, int dig)
{
    return pgh[index][dig];
}
```

```
leaq (%rdi,%rdi,4), %rax # 5*index
addl %rax, %rsi # 5*index+dig
movl pgh(,%rsi,4), %eax # M[pgh + 4*(5*index+dig)]
```

Array Elements

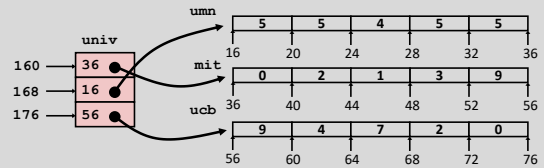
- `pgh[index][dig]` is `int`
- Address: `pgh + 20*index + 4*dig`
 $\ast = \text{pgh} + 4 \ast (5 \ast \text{index} + \text{dig})$

Multi-Level Array Example

```
zip_dig umn = { 5, 5, 4, 5, 5 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

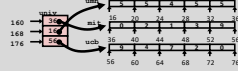
```
#define UCOUNT 3
int *univ[UCOUNT] = {mit, umn, ucb};
```

- Variable `univ` denotes array of 3 elements
- Each element is a pointer
 - 8 bytes
- Each pointer points to array of `int`'s



Element Access in Multi-Level Array

```
int get_univ_digit
(size_t index, size_t digit)
{
    return univ[index][digit];
}
```



```
salq $2, %rsi # 4*digit
addq univ(,%rdi,8), %rsi # p = univ[index] + 4*digit
movl (%rsi), %eax # return *p
ret
```

Computation

- Element access `Mem[Mem[univ+8*index]+4*digit]`
- Must do two memory reads
 - First get pointer to row array
 - Then access element within array

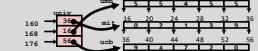
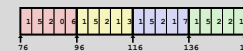
Array Element Accesses

Nested array

```
int get_pgh_digit
(size_t index, size_t digit)
{
    return pgh[index][digit];
}
```

Multi-level array

```
int get_univ_digit
(size_t index, size_t digit)
{
    return univ[index][digit];
}
```



Accesses looks similar in C, but address computations very different:

`Mem[pgh+20*index+4*digit]` `Mem[Mem[univ+8*index]+4*digit]`

N X N Matrix Code

Fixed dimensions

- Know value of `N` at compile time

Variable dimensions, explicit indexing

- Traditional way to implement dynamic arrays

Variable dimensions, implicit indexing

- Now supported by gcc

```
#define N 16
typedef int fix_matrix[N][N];
/* Get element a[i][j] */
int fix_ele(fix_matrix a,
            size_t i, size_t j)
{
    return a[i][j];
}
```

```
#define IDX(n, i, j) ((i)*(n)+(j))
/* Get element a[i][j] */
int vec_ele(size_t n, int *a,
            size_t i, size_t j)
{
    return a[IDX(n,i,j)];
}
```

```
/* Get element a[i][j] */
int var_ele(size_t n, int a[n][n],
            size_t i, size_t j) {
    return a[i][j];
}
```

16 X 16 Matrix Access

Array Elements

- Address $A + i \ast (C \ast K) + j \ast K$
- $C = 16, K = 4$

```
/* Get element a[i][j] */
int fix_ele(fix_matrix a, size_t i, size_t j) {
    return a[i][j];
}
```

```
# a in %rdi, i in %rsi, j in %rdx
salq $6, %rsi # 64*i
addq %rsi, %rdi # a + 64*i
movl (%rdi,%rdx,4), %eax # M[a + 64*i + 4*j]
ret
```

n X n Matrix Access

■ Array Elements

- Address $A + i * (C * K) + j * K$
- $C = n, K = 4$
- Must perform integer multiplication

```
/* Get element a[i][j] */
int var_ele(size_t n, int a[n][n], size_t i, size_t j)
{
    return a[i][j];
}
```

```
# n in %rdi, a in %rsi, i in %rdx, j in %rcx
imulq    %rdx, %rdi    # n*i
leaq     (%rsi,%rdi,4), %rax # a + 4*n*i
movl     (%rax,%rcx,4), %eax # a + 4*n*i + 4*j
ret
```

Today

■ Arrays

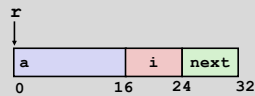
- One-dimensional
- Multi-dimensional (nested)
- Multi-level

■ Structures

- Allocation
- Access
- Alignment

Structure Representation

```
struct rec {
    int a[4];
    size_t i;
    struct rec *next;
};
```



■ Structure represented as block of memory

- Big enough to hold all of the fields

■ Fields ordered according to declaration

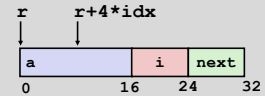
- Even if another ordering could yield a more compact representation

■ Compiler determines overall size + positions of fields

- Machine-level program has no understanding of the structures in the source code

Generating Pointer to Structure Member

```
struct rec {
    int a[4];
    size_t i;
    struct rec *next;
};
```



■ Generating Pointer to Array Element

- Offset of each structure member determined at compile time
- Compute as $r + 4 * idx$

```
int *get_ap
(struct rec *r, size_t idx)
{
    return &r->a[idx];
}
```

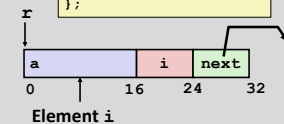
```
# r in %rdi, idx in %rsi
leaq    (%rdi,%rsi,4), %rax
ret
```

Following Linked List

■ C Code

```
void set_val
(struct rec *r, int val)
{
    while (r) {
        int i = r->i;
        r->a[i] = val;
        r = r->next;
    }
}
```

```
struct rec {
    int a[4];
    int i;
    struct rec *next;
};
```

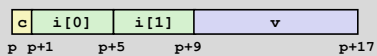


Register	Value
%rdi	r
%rsi	val

```
.L11:
movslq    16(%rdi), %rax    # loop:
# i = M[r+16]
movl      %esi, (%rdi,%rax,4) # M[r+4*i] = val
movq      24(%rdi), %rdi    # r = M[r+24]
testq     %rdi, %rdi        # Test r
jne        .L11             # if !=0 goto loop
```

Structures & Alignment

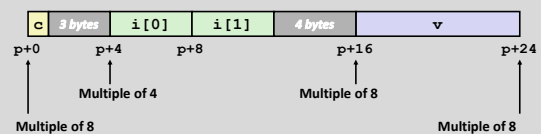
■ Unaligned Data



```
struct S1 {
    char c;
    int i[2];
    double v;
} *p;
```

■ Aligned Data

- Primitive data type requires K bytes
- Address must be multiple of K



Alignment Principles

- **Aligned Data**
 - Primitive data type requires K bytes
 - Address must be multiple of K
 - Required on some machines; advised on x86-64
- **Motivation for Aligning Data**
 - Memory accessed by (aligned) chunks of 4 or 8 bytes (system dependent)
 - Inefficient to load or store datum that spans quad word boundaries
 - Virtual memory trickier when datum spans 2 pages
- **Compiler**
 - Inserts gaps in structure to ensure correct alignment of fields

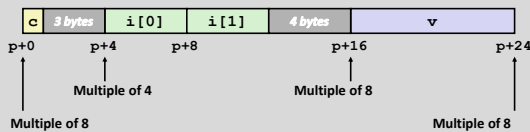
Specific Cases of Alignment (x86-64)

- **1 byte:** `char`, ...
 - no restrictions on address
- **2 bytes:** `short`, ...
 - lowest 1 bit of address must be 0_2
- **4 bytes:** `int`, `float`, ...
 - lowest 2 bits of address must be 00_2
- **8 bytes:** `double`, `long`, `char *`, ...
 - lowest 3 bits of address must be 000_2
- **16 bytes:** `long double` (GCC on Linux)
 - lowest 4 bits of address must be 0000_2

Satisfying Alignment with Structures

- **Within structure:**
 - Must satisfy each element's alignment requirement
- **Overall structure placement**
 - Each structure has alignment requirement K
 - K = Largest alignment of any element
 - Initial address & structure length must be multiples of K
- **Example:**
 - $K = 8$, due to `double` element

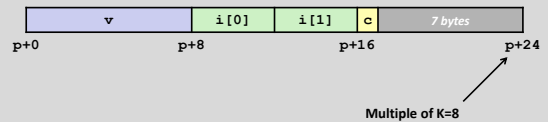
```
struct S1 {
    char c;
    int i[2];
    double v;
} *p;
```



Meeting Overall Alignment Requirement

- For largest alignment requirement K
- Overall structure must be multiple of K

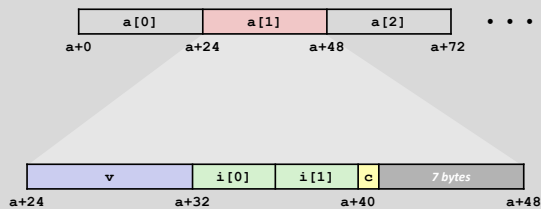
```
struct S2 {
    double v;
    int i[2];
    char c;
} *p;
```



Arrays of Structures

- Overall structure length multiple of K
- Satisfy alignment requirement for every element

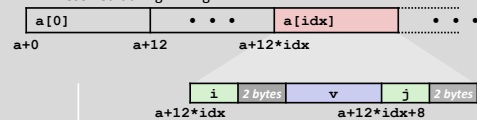
```
struct S2 {
    double v;
    int i[2];
    char c;
} a[10];
```



Accessing Array Elements

- Compute array offset $12 \cdot \text{idx}$
 - `sizeof(S3)`, including alignment spacers
- Element j is at offset 8 within structure
- Assembler gives offset $a+8$
 - Resolved during linking

```
struct S3 {
    short i;
    float v;
    short j;
} a[10];
```



```
short get_j(int idx)
{
    return a[idx].j;
}
```

```
# %rdi = idx
leaq (%rdi,%rdi,2),%rax # 3*idx
movzwl a+8(,%rax,4),%eax
```

Saving Space

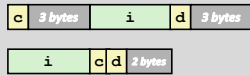
- Put large data types first

```
struct S4 {  
  char c;  
  int i;  
  char d;  
} *p;
```



```
struct S5 {  
  int i;  
  char c;  
  char d;  
} *p;
```

- Effect (K=4)



Summary

- Arrays
 - Elements packed into contiguous region of memory
 - Use index arithmetic to locate individual elements
- Structures
 - Elements packed into single region of memory
 - Access using offsets determined by compiler
 - Possible require internal and external padding to ensure alignment
- Combinations
 - Can nest structure and array code arbitrarily