CSCI 5105

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Today

- Communication in Distributed Systems
 Overview
 - Types
- Remote Procedure Calls (RPC)
- Remote Method Invocation (RMI)

Communication

How do program modules/processes communicate on a single machine?

Communication in Distributed Systems

- "Distributed" processes
- Located on different machines
- Need communication mechanisms
- Goal: Hide distributed nature as far as possible

Communication in Distributed Systems

- Networking primitives and protocols (e.g.: TCP/IP)
- Advanced communication models: Built on networking primitives
 - Remote Procedure Calls (RPC)
 - Remote Method Invocation (RMI)
 - Messages
 - Multicast

Types of Communication

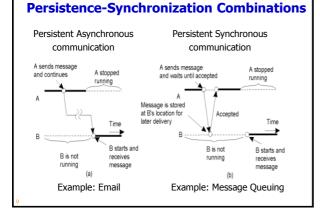
- Defined by two main properties:
 - Persistence
 - Synchronization

Persistence

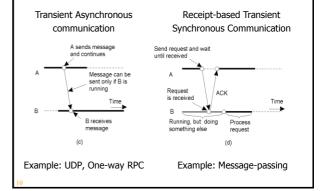
- Persistent communication
 - Messages are stored until receiver is ready
 - Sender/receiver don't have to be up at the same time
- Transient communication
 - Message is stored only so long as both sending/receiving applications are executing
 - Discard message if it can't be delivered to receiver

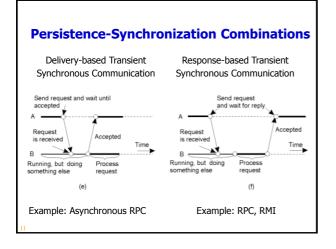
Synchronization

- Synchronous communication
 - Sender blocks until message is delivered to receiver
 - Variant: block until receiver processes the message
- Asynchronous communication
 - Sender continues immediately after it has submitted the message
- Several combinations of persistence and synchronization



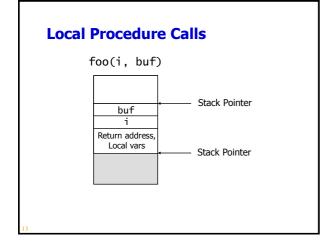
Persistence-Synchronization Combinations





Remote Procedure Calls (RPC)

- Goal: Make distributed computation look like centralized computation
- Idea: Allow processes to call procedures on other machines
 - Make it appear like normal procedure calls



RPC Operation

Challenges:

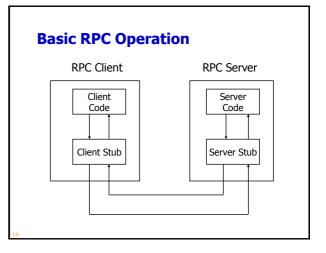
- Hide details of communication
- Pass parameters transparently

Stubs

- Hide communication details
- Client and server stubs
- Marshalling
 - Flattening and parameter passing

Stubs

- Code that communicates with the remote side
- Client stub:
 - Converts function call to remote communication
 - Passes parameters to server machine
 - Receives results
- Server stub:
 - Receives parameters and request from client
 - Calls the desired server function
 - Returns results to client



Parameter Passing: Local Procedures

- Pass-by-value
 - Original variable is not modified
 - E.g.: integers, chars
- Pass-by-reference
 - Passing a pointer
 - Value may be changed
 - E.g.: Arrays
- Pass-by-copy/restore
 - Copy is modified and overwritten to the original
 - E.g.: in-out parameters in Ada

Parameter Passing: RPC

- Pass-by-value
 - Send the value in standard format
- Pass-by-reference
 - Can we pass pointers?
- What about complex data structures (linked lists, trees, graphs)?

Marshalling

- Converting parameters into a byte stream
- Problems:
 - Heterogeneous data formats: Big-endian vs. little-endian
 - Type of parameter passing: by-value vs. by-reference

Heterogeneous Data Formats

- Use a standard data format
- Examples: Network byte order, XDR (Extended Data Representation)
- Decide on a protocol for parameter ordering

Stub Generation

Most stubs are similar in functionality

- Handle communication and marshalling
- Differences are in the main server-client code
- Application needs to know only stub interface
- Interface Definition Language (IDL)
 - Allows interface specification
 - IDL compiler generates the stubs automatically

Binding

- How does the client stub find the server stub?
 Needs to know remote IP address/port no.
- Port mapper
 - Daemon on server machine maintaining server bindings
 - Listens on a well-known port
- Server stub registers its port no. and service name with portmapper
 - Client gets this binding by querying portmapper
- Server's IP address can be obtained from a directory service

Synchronous RPC

- RPC Performed in a synchronous manner
 - Client blocks until results come back
- What if client wants to do something else?

RPC Variants

- Asynchronous RPC
 - Server sends ACK as soon as request is receivedExecutes procedure later
- Deferred synchronous RPC
 - Use two asynchronous RPCs
 - Server sends reply via second asynchronous RPC
 - Callback to the client
- One-way RPC
 - Client does not even wait for an ACK from the server

Multicast RPC

- RPC sent to multiple servers
 - Could use multiple concurrent one-way or asynchronous RPCs
- Why use multicast RPC?
- How to handle responses?
 - First one or majority response
 - Aggregate multiple responses

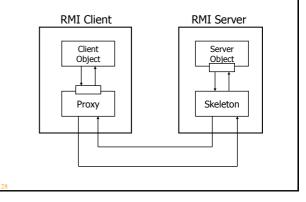
Remote Method Invocation (RMI)

- RPCs applied to distributed objects
- Class: object-oriented abstraction
- Object: instance of class
 - Encapsulates data
 - Exports methods: operations on data
 - Separation between interface and implementation

Distributed Objects

- Object resides on one machine, interfaces reside on other machines
- Remote object: State on different machine
- Local object: State on the same machine
- RMIs allow invoking methods of remote objects
 - Use proxies, skeletons, binding
 - Allow passing of object references as parameters

Basic RMI Operation



Proxies and Skeletons

- Proxy: client stub
 - Maintains server ID, endpoint, object ID
 - Does parameter marshalling
 - In practice, can be downloaded/constructed on the fly
- Skeleton: server stub
 - Does demarshalling and passes parameters to server
 - Sends result to proxy

Binding a Client to an Object

- Loading a proxy in client address space
- Implicit binding:
 - Bound automatically on object reference resolution
- Explicit binding:
 - Client has to first bind object
 - Call method after binding

