CSci 5271
Introduction to Computer Security
Networking (cont'd) and cryptography

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Preview question

Which of these is a cryptographic primitive based on a Feistel cipher design?

A. DES
B. AES
C. DSA
D. CBC
E. HMAC

Outline

TCP

- Transmission Control Protocol: provides reliable bidirectional stream abstraction
- Packets have sequence numbers, acknowledged in order
- Missed packets resent later

Flow and congestion control

- Flow control: match speed to slowest link
  - "Window" limits number of packets sent but not ACKed
- Congestion control: avoid traffic jams
  - Lost packets signal congestion
  - Additive increase, multiplicative decrease of rate

Routing

- Where do I send this packet next?
  - Table from address ranges to next hops
- Core Internet routers need big tables
- Maintained by complex, insecure, cooperative protocols
  - Internet-level algorithm: BGP (Border Gateway Protocol)

Below IP: ARP

- Address Resolution Protocol maps IP addresses to lower-level address
  - E.g., 48-bit Ethernet MAC address
- Based on local-network broadcast packets
- Complex Ethernets also need their own routing (but called switches)

DNS

- Domain Name System: map more memorable and stable string names to IP addresses
- Hierarchically administered namespace
  - Like Unix paths, but backwards
- `.edu` server delegates to `.umn.edu` server, etc.

DNS caching and reverse DNS

- To be practical, DNS requires caching
  - Of positive and negative results
- But, cache lifetime limited for freshness
- Also, reverse IP to name mapping
  - Based on special top-level domain, IP address written backwards

Classic application: remote login

- Killer app of early Internet: access supercomputers at another university
- Telnet: works cross-OS
  - Send character stream, run regular login program
- rlogin: BSD Unix
  - Can authenticate based on trusting computer connection comes from
  - (Also rsh, rcp)

Outline

Packet sniffing

- Watch other people's traffic as it goes by on network
- Easiest on:
  - Old-style broadcast (thin, "hub") Ethernet
  - Wireless
- Or if you own the router

Forging packet sources

- Source IP address not involved in routing, often not checked
- Change it to something else!
- Might already be enough to fool a naive UDP protocol

TCP spoofing

- Forging source address only lets you talk, not listen
- Old attack: wait until connection established, then DoS one participant and send packets in their place
- Frustrated by making TCP initial sequence numbers unpredictable
  - But see Oakland'12, WOOT'12 for fancier attacks, keyword "off-path"

ARP spoofing

- Impersonate other hosts on local network level
- Typical ARP implementations stateless, don't mind changes
- Now you get victim's traffic, can read, modify, resend

rlogin and reverse DNS

- rlogin uses reverse DNS to see if originating host is on whitelist
- How can you attack this mechanism with an honest source IP address?

rlogin and reverse DNS

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- How can you attack this mechanism with an honest source IP address?
- Remember, ownership of reverse-DNS is by IP address

Outline

Midterms

- Graded midterms will be given back on Monday
- Grades will also be on Canvas by then
- There may be a difficulty adjustment

Project meetings

- Next round of meetings next week 10/28-11/1
- Mostly same times as before, will confirm by email

Outline

-ography, -ology, -analysis

- Cryptography (narrow sense): designing encryption
- Cryptanalysis: breaking encryption
- Cryptology: both of the above
- Code (narrow sense): word-for-concept substitution
- Cipher: the "codes" we actually care about

Caesar cipher

- Advance three letters in alphabet: $A \to D, B \to E, \ldots$
- Decrypt by going back three letters
- Internet-era variant: rot-13
- Easy to break if you know the principle

Keys and Kerckhoffs's principle

- The only secret part of the cipher is a *key*
- Security does not depend on anything else being secret
- Modern (esp. civilian, academic) crypto embraces openness quite strongly

Symmetric vs. public key

- Symmetric key (today's lecture): one key used by all participants
- Public key: one key kept secret, another published
  - Techniques invented in 1970s
  - Makes key distribution easier
  - Depends on fancier math

Goal: secure channel

- Leaks no content information
  - Not protected: size, timing
- Messages delivered intact and in order
  - Or not at all
- Even if an adversary can read, insert, and delete traffic

One-time pad

- Secret key is truly random data as long as message
- Encrypt by XOR (more generally addition mod alphabet size)
- Provides perfect, "information-theoretic" secrecy
- No way to get around key size requirement

Computational security

- More realistic: assume adversary has a limit on computing power
- Secure if breaking encryption is computationally infeasible
  - E.g., exponential-time brute-force search
- Ties cryptography to complexity theory

Key sizes and security levels

- Difficulty measured in powers of two, ignore small constant factors
- Power of attack measured by number of steps, aim for better than brute force
- 2^{32} definitely too easy, probably 2^{64} too
- Modern symmetric key size: at least 2^{128}

Crypto primitives

- Base complicated systems on a minimal number of simple operations
- Designed to be fast, secure in wide variety of uses
- Study those primitives very intensely

Attacks on encryption

- Known ciphertext
  - Weakest attack
- Known plaintext (and corresponding ciphertext)
- Chosen plaintext
- Chosen ciphertext (and plaintext)
  - Strongest version: adaptive

Certificational attacks

- Good primitive claims no attack more effective than brute force
- Any break is news, even if it's not yet practical
  - Canary in the coal mine
- E.g., $2^{126.1}$ attack against AES-128
- Also watched: attacks against simplified variants

Fundamental ignorance

- We don't really know that any computational cryptosystem is secure
- Security proof would be tantamount to proving $P \neq NP$
- Crypto is fundamentally more uncertain than other parts of security

Relative proofs

- Prove security under an unproved assumption
- In symmetric crypto, prove a construction is secure if the primitive is
  - Often the proof looks like: if the construction is insecure, so is the primitive
- Can also prove immunity against a particular kind of attack

Random oracle paradigm

- Assume ideal model of primitives: functions selected uniformly from a large space
  - Anderson: elves in boxes
- Not theoretically sound; assumption cannot be satisfied
- But seems to be safe in practice

Pseudorandomness and distinguishers

- Claim: primitive cannot be distinguished from a truly random counterpart
  - In polynomial time with non-negligible probability
- We can build a distinguisher algorithm to exploit any weakness
- Slightly too strong for most practical primitives, but a good goal

Open standards

- How can we get good primitives?
- Open-world best practice: run competition, invite experts to propose then attack
- Run by neutral experts, e.g. US NIST
- Recent good examples: AES, SHA-3

A certain three-letter agency

- National Security Agency (NSA): has primary responsibility for "signals intelligence"
- Dual-mission tension:
  - Break the encryption of everyone in the world
  - Help US encryption not be broken by foreign powers

Outline

Stream ciphers

- Closest computational version of one-time pad
- Key (or seed) used to generate a long pseudorandom bitstream
- Closely related: cryptographic RNG

Shift register stream ciphers

- Linear-feedback shift register (LFSR): easy way to generate long pseudorandom sequence
  - But linearity allows for attack
- Several ways to add non-linearity
- Common in constrained hardware, poor security record

RC4

- Fast, simple, widely used software stream cipher
  - Previously a trade secret, also "ARCFOUR"
- Many attacks, none yet fatal to careful users (e.g. TLS)
  - Famous non-careful user: WEP
- Now deprecated, not recommended for new uses

Encryption $\neq$ integrity

- Encryption protects secrecy, not message integrity
- For constant-size encryption, changing the ciphertext just creates a different plaintext
- How will your system handle that?
- Always need to take care of integrity separately

Stream cipher mutability

- Strong example of encryption vs. integrity
- In stream cipher, flipping a ciphertext bit flips the corresponding plaintext bit, only
- Very convenient for targeted changes

Stream cipher assessment

- Currently out of fashion as a primitive in software
- Not inherently insecure
  - Other common pitfall: must not reuse key(stream)
- Currently no widely vetted primitives

Outline

Basic idea

- Encryption/decryption for a fixed sized block
- Insecure if block size is too small
  - Barely enough: 64 bits; current standard: 128
- Reversible, so must be one-to-one and onto function

Pseudorandom permutation

- Ideal model: key selects a random invertible function
- I.e., permutation (PRP) on block space
  - Note: not permutation on bits
- "Strong" PRP: distinguisher can decrypt as well as encrypt

Confusion and diffusion

- Basic design principles articulated by Shannon
- Confusion: combine elements so none can be analyzed individually
- Diffusion: spread the effect of one symbol around to others
- Iterate multiple *rounds* of transformation

Substitution/permutation network

- Parallel structure combining reversible elements:
- Substitution: invertible lookup table ("S-box")
- Permutation: shuffle bits

AES

- Advanced Encryption Standard: NIST contest 2001
  - Developed under the name Rijndael
- 128-bit block, 128/192/256-bit key
- Fast software implementation with lookup tables (or dedicated insns)
- Allowed by US government up to Top Secret

Feistel cipher

- Split block in half, operate in turn: $(L_{i+1}, R_{i+1}) = (R_i, L_i \oplus F(R_i, K_i))$
- Key advantage: F need not be invertible
  - Also saves space in hardware
- Luby-Rackoff: if F is pseudo-random, 4 or more rounds gives a strong PRP

DES

- Data Encryption Standard: AES predecessor 1977-2005
- 64-bit block, 56-bit key
- Implementable in 70s hardware, not terribly fast in software
- Triple DES variant still used in places

Some DES history

- Developed primarily at IBM, based on an earlier cipher named "Lucifer"
- Final spec helped and "helped" by the NSA
  - Argued for smaller key size
  - S-boxes tweaked to avoid a then-secret attack
- Eventually victim to brute-force attack

DES brute force history

1977	est. $20m cost custom hardware
1993	est. $1m cost custom hardware
1997	distributed software break
1998	$250k built ASIC hardware
2006	$10k FPGAs
2012	as-a-service against MS-CHAPv2

Double encryption?

- Combine two different block ciphers?
  - Belt and suspenders
- Anderson: don't do it
- FS&K: could do it, not a recommendation
- Maurer and Massey (J.Crypt'93): might only be as strong as first cipher

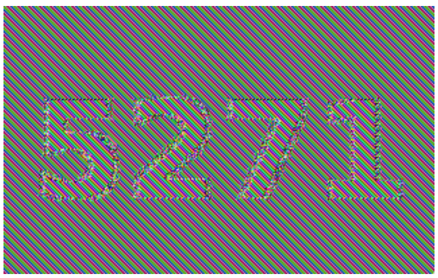
Modes of operation

- How to build a cipher for arbitrary-length data from a block cipher
- Many approaches considered
  - For some reason, most have three-letter acronyms
- More recently: properties susceptible to relative proof

ECB

- Electronic CodeBook
- Split into blocks, apply cipher to each one individually
- Leaks equalities between plaintext blocks
- Almost never suitable for general use

Do not use ECB



CBC

- Cipher Block Chaining
- $C_i = E_K(P_i \oplus C_{i-1})$
- Probably most popular in current systems
- Plaintext changes propagate forever, ciphertext changes only one block

CBC: getting an IV

- C_0 is called the initialization vector (IV)
  - Must be known for decryption
- IV should be random-looking
  - To prevent first-block equalities from leaking (lesser version of ECB problem)
- Common approaches
  - Generate at random
  - Encrypt a nonce

Stream modes: OFB, CTR

- Output FeedBack: produce keystream by repeatedly encrypting the IV
  - Danger: collisions lead to repeated keystream
- Counter: produce from encryptions of an incrementing value
  - Recently becoming more popular: allows parallelization and random access

Outline

Brief introduction to networking, (cont'd)
Some classic network attacks
Announcements intermission
Crypto basics
Stream ciphers
Block ciphers and modes of operation
Hash functions and MACs
Building a secure channel

Ideal model

- Ideal crypto hash function: pseudorandom function
  - Arbitrary input, fixed-size output
- Simplest kind of elf in box, theoretically very convenient
- But large gap with real systems: better practice is to target particular properties

Kinds of attacks

- Pre-image, "inversion": given y, find x such that $H(x) = y$
- Second preimage, targeted collision: given x, $H(x)$, find $x' \neq x$ such that $H(x') = H(x)$
- (Free) collision: find x_1, x_2 such that $H(x_1) = H(x_2)$

Birthday paradox and attack

- There are almost certainly two people in this classroom with the same birthday
- n people have $\binom{n}{2} = \Theta(n^2)$ pairs
- So only about $\sqrt{n}$ expected for collision
- "Birthday attack" finds collisions in any function

Security levels

- For function with k-bit output:
- Preimage and second preimage should have complexity 2^k
- Collision has complexity $2^{k/2}$
- Conservative: use hash function twice as big as block cipher key
  - Though if you're paranoid, cipher blocks can repeat too

Non-cryptographic hash functions

- The ones you probably use for hash tables
- CRCs, checksums
- Output too small, but also not resistant to attack
- E.g., CRC is linear and algebraically nice

Short hash function history

- On the way out: MD5 (128 bit)
  - Flaws known, collision-finding now routine
- SHA(-0): first from NIST/NSA, quickly withdrawn
  - Likely flaw discovered 3 years later
- SHA-1: fixed SHA-0, 160-bit output.
- 2^{60} collision attack described in 2013
  - First public collision found (using 6.5 kCPU yr) in 2017

Length extension problem

- MD5, SHA1, etc., computed left to right over blocks
- Can sometimes compute $H(a \parallel b)$ in terms of $H(a)$
  - $\parallel$ means bit string concatenation
- Makes many PRF-style constructions insecure

SHA-2 and SHA-3

- SHA-2: evolutionary, larger, improvement of SHA-1
  - Exists as SHA-$\{224, 256, 384, 512\}$
  - But still has length-extension problem
- SHA-3: chosen recently in open competition like AES
  - Formerly known as Keccak, official standard Aug. 2015
  - New design, fixes length extension
  - Not yet very widely used

MAC: basic idea

- Message authentication code: similar to hash function, but with a key
- Adversary without key cannot forge MACs
- Strong definition: adversary cannot forge anything, even given chosen-message MACs on other messages

CBC-MAC construction

- Same process as CBC encryption, but:
  - Start with IV of 0
  - Return only the last ciphertext block
- Both these conditions needed for security
- For fixed-length messages (only), as secure as the block cipher

HMAC construction

- $H(K \parallel M)$: insecure due to length extension
  - Still not recommended: $H(M \parallel K)$, $H(K \parallel M \parallel K)$
- HMAC: $H(K \oplus a \parallel H(K \oplus b \parallel M))$
- Standard $a = \mathtt{0x5c}^*$, $b = \mathtt{0x36}^*$
- Probably the most widely used MAC

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Session keys

- Don't use your long term password, etc., directly as a key
- Instead, *session key* used for just one channel
- In modern practice, usually obtained with public-key crypto
- Separate keys for encryption and MACing

Order of operations

- Encrypt and MAC ("in parallel")
  - Safe only under extra assumptions on the MAC
- Encrypt then MAC
  - Has cleanest formal safety proof
- MAC then Encrypt
  - Preferred by FS&K for some practical reasons
  - Can also be secure

Authenticated encryption modes

- Encrypting and MACing as separate steps is about twice as expensive as just encrypting
- "Authenticated encryption" modes do both at once
  - Newer (circa 2000) innovation, many variants
- NIST-standardized and unpatented: Galois Counter Mode (GCM)

Ordering and message numbers

- Also don't want attacker to be able to replay or reorder messages
- Simple approach: prefix each message with counter
- Discard duplicate/out-of-order messages

Padding

- Adjust message size to match multiple of block size
- To be reversible, must sometimes make message longer
- E.g.: for 16-byte block, append either 1, or $2\ 2$, or $3\ 3\ 3$, up to 16 "16" bytes

Padding oracle attack

- Have to be careful that decoding of padding does not leak information
- E.g., spend same amount of time MACing and checking padding whether or not padding is right
- Remote timing attack against CBC TLS published 2013

Don’t actually reinvent the wheel

- This is all implemented carefully in OpenSSL, SSH, etc.
- Good to understand it, but rarely sensible to reimplement it
- You’ll probably miss at least one of decades’ worth of attacks

Next time

- Public-key encryption protocols
- More about provable security and appropriate paranoia