### CSci 5271 Introduction to Computer Security Day 12: OS security: higher assurance

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### Outline

OS trust and assurance

Announcements intermission

### Trusted and trustworthy

- Part of your system is trusted if its failure can break your security
- Thus, OS is almost always trusted
- Real question: is it trustworthy?
- Distinction not universally observed: trusted boot, Trusted Solaris, etc.

### Trusted (I/O) path

- How do you know you're talking to the right software?
- And no one is sniffing the data?
- 🖲 Example: Trojan login screen
  - Or worse: unlock screensaver with root password
  - Origin of "Press Ctrl-Alt-Del to log in"

### Minimizing trust

**E** Kernel  $\rightarrow$  microkernel  $\rightarrow$  nanokernel

- Reference monitor concept
- TCB size: measured relative to a policy goal

**O** Reference monitor  $\subseteq$  TCB

But hard to build monitor for all goals

### How to gain assurance

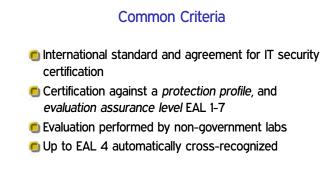
- Use for a long time
- 🖲 Testing
- 🖲 Code / design review
- Third-party certification
- Formal methods / proof

# Evaluation / certification

- Testing and review performed by an independent party
- Goal: separate incentives, separate accountability
- Compare with financial auditing
- Watch out for: form over substance, misplaced incentives

# Orange book OS evaluation Trusted Computer System Evaluation Criteria Minimal protection Discretionary protection C2 adds, e.g., secure audit over C1

- B. Mandatory protection
  - B1<B2<B3: stricter classic MLS</p>
- A. Verified protection



### Common Criteria, Anderson's view

Many profiles don't specify the right things
 OSes evaluated only in unrealistic environments
 E.g., unpatched Windows XP with no network attacks

- Corruption, Manipulation, and Inertia"
  - Pernicious innovation: evaluation paid for by vendor
     Labs beholden to national security apparatus

### Formal methods and proof

- Can math come to the rescue?
- Checking design vs. implementation
- Automation possible only with other tradeoffs E.g., bounded size model
- Starting to become possible: machine-checked proof

### Proof and complexity

- Formal proof is only feasible for programs that are small and elegant
- If you honestly care about assurance, you want your TCB small and elegant anyway
- Should provability further guide design?

### Some hopeful proof results

seL4 microkernel (SOSP'09 and ongoing)
 7.5 kL C, 200 kL proof, 160 bugs fixed, 25 person years
 CompCert C-subset compiler (PLDI'06 and ongoing)
 RockSalt SFI verifier (PLDI'12)

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### Note to early readers

- This is the section of the slides most likely to change in the final version
- If class has already happened, make sure you have the latest slides for announcements
- In particular, the BCMTA vulnerability announcement is embargoed